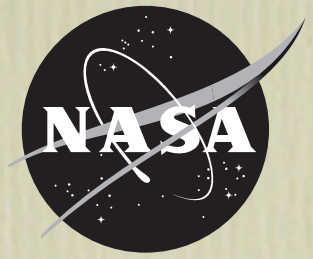


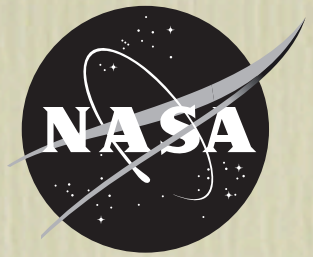
Kerberized Certificate Issuance Protocol (KX509)

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Jet Propulsion Laboratory

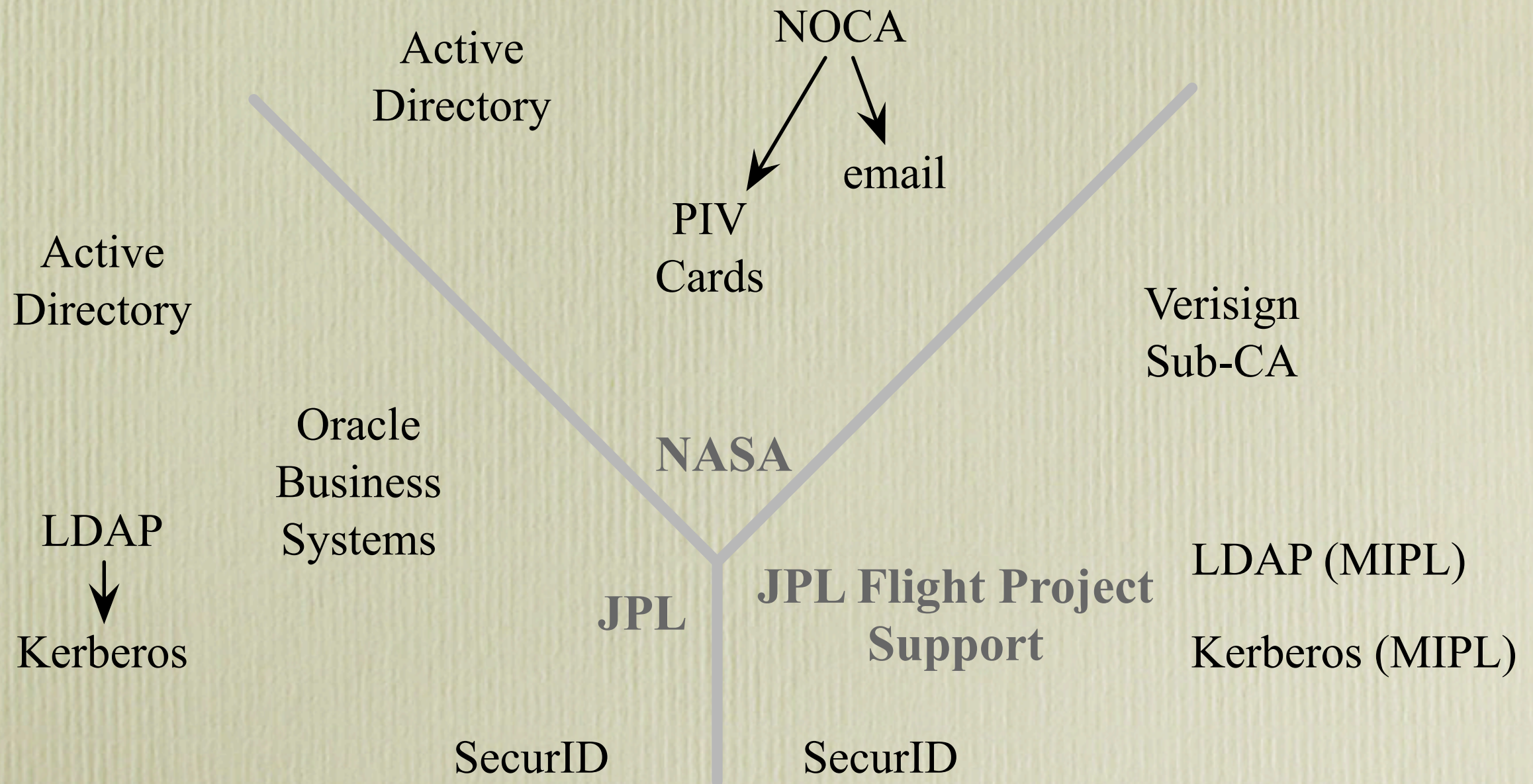


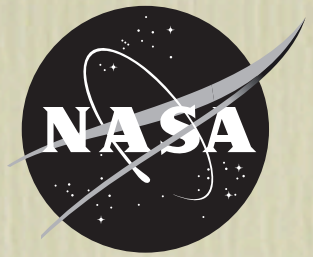
Overview and Purpose

- KX509 is a wire protocol for using Kerberos tickets to acquire X.509 certificates.
 - Kind of the opposite of PKINIT
- Where both X.509 and Kerberos are in use, want to guarantee they both authoritatively refer to the same entities.
- Already in use at several large institutions.



Identity Infrastructures at JPL





Protocol Overview

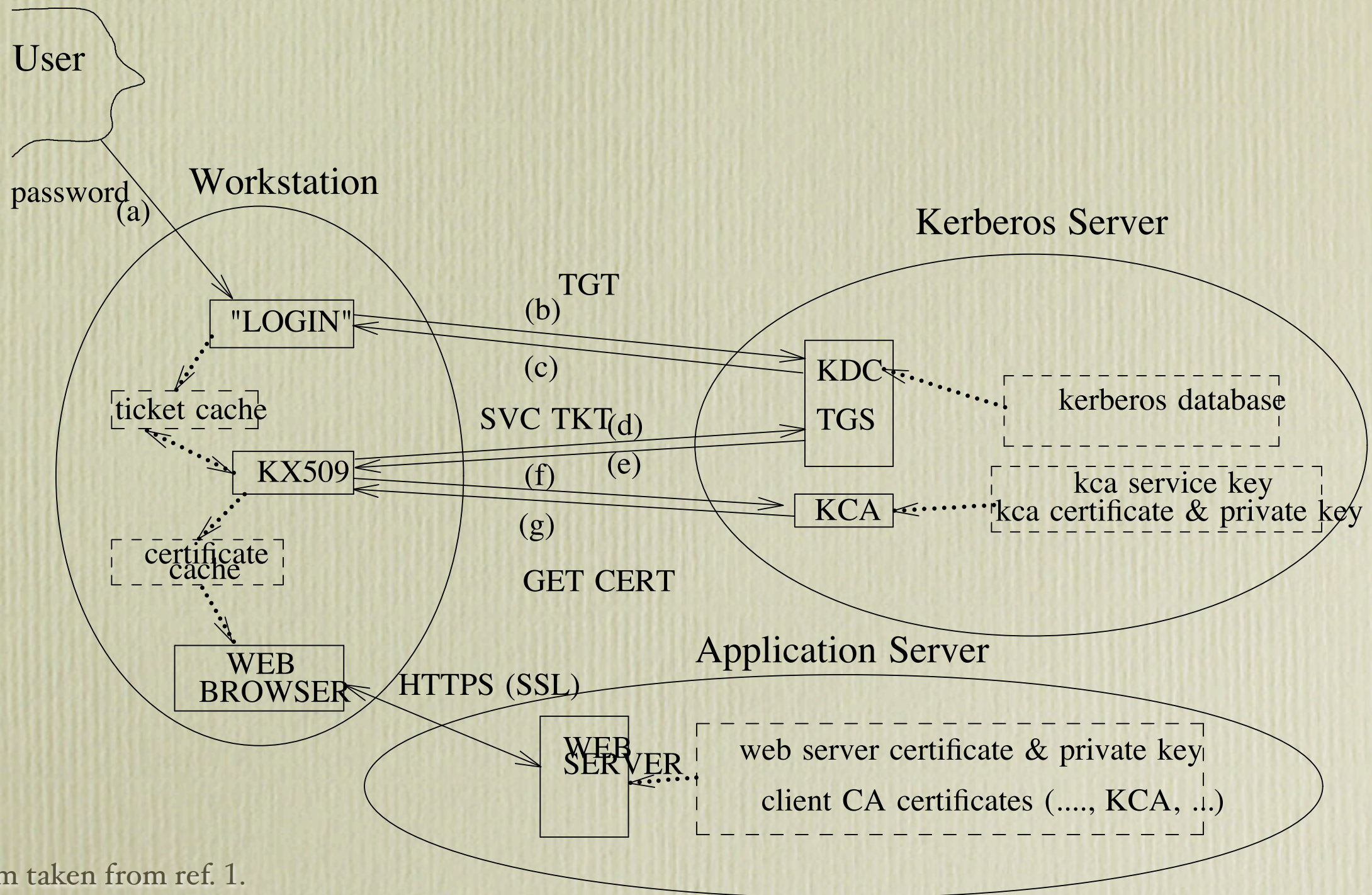
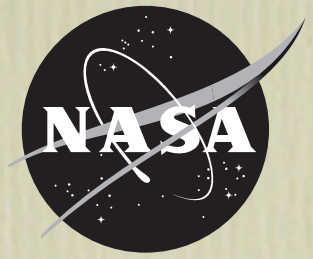
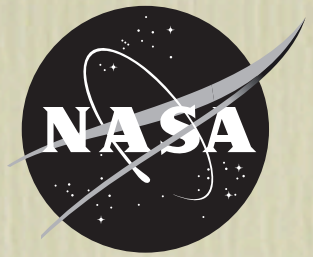


Diagram taken from ref. 1.



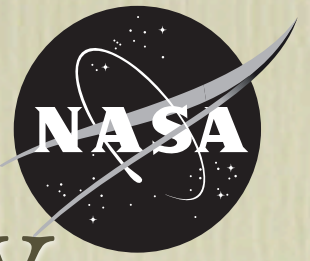
Client Implementation

- Command line utilities
 - `kx509`
 - Generates pub/private keys
 - Does protocol exchange with KCA
 - Stores certificate in Kerberos credential cache
 - `kxlist`
 - List certificates stored in the Kerberos credential cache
- PKCS-11 library
 - Implements PKI support using the cert/key in the credential cache.
- Interest in having the library get the cert when opened.



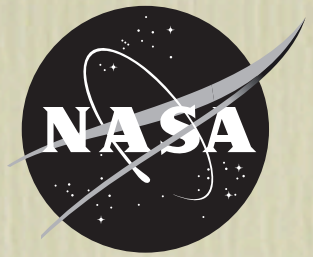
Protocol Description – Request

- UDP packet containing
 - version string (2.0)
 - ap-req – normal Kerberos stuff.
 - pk-hash – HMAC/SHA-1 of the version string and the pk-key.
 - pk-key – RSA public key.
 - UMICH implementation supports DSA keys, but not used “in the wild”.
- Nothing is encrypted



Protocol Description – Reply

- UDP packet containing
 - version string (2.0)
 - error-code – 0 (absent) means OK.
 - hash – HMAC/SHA-1 of the reply fields present.
 - certificate – X.509 certificate.
 - e-text – error message.

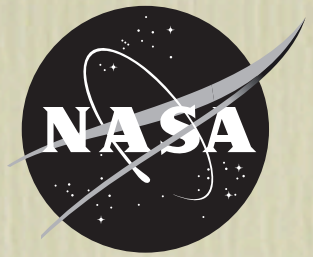


Reply Constraints

- All fields are nominally optional, but only the following combinations are allowed:

certificate			hash
	error-code	e-text	hash
	error-code	e-text	

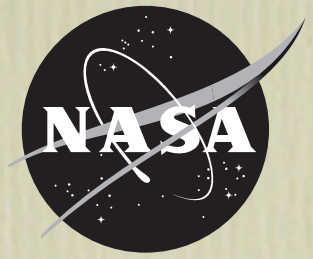
- The certificate should contain:
 - Subject name unique to requestor.
 - Unique serial number (across all KCA's).
 - An extension identifying the original Kerberos identity
 - `id-pkinit-san` – preferred
 - `kcaAuthRealm` – realm only
 - `userPrincipalName` – similar to `id-pkinit-san`.



Observed Deployment

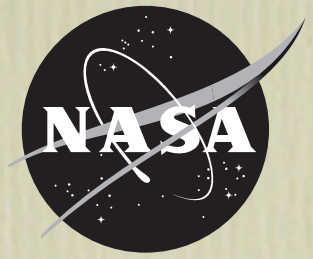
```
Version: 3 (0x2)
Serial Number: 30358893 (0x1cf3d6d)
Signature Algorithm: sha1WithRSAEncryption
Issuer: DC=gov, DC=fnal, O=Fermilab, OU=Certificate Authorities, CN=Kerberized CA HSM
Validity
  Not Before: Sep 23 18:48:37 2010 GMT
  Not After : Oct  1 15:10:31 2010 GMT
Subject: DC=gov, DC=fnal, O=Fermilab, OU=People, CN=Matt Crawford, CN=UID:crawdad
Subject Public Key Info:
  Public Key Algorithm: rsaEncryption
  RSA Public Key: (1024 bit)
    Modulus (1024 bit): . . .
    Exponent: 65537 (0x10001)
X509v3 extensions:
  X509v3 Basic Constraints: critical
    CA:FALSE
  X509v3 Key Usage: critical
    Digital Signature, Key Encipherment
  Netscape Cert Type:
    SSL Client
  Netscape Comment:
    Certificate issued by Fermilab KCA
  X509v3 Issuer Alternative Name:
    email:nightwatch@fnal.gov
  Netscape CA Policy Url:
    URL:http://security.fnal.gov/policies/pki\_policy\_certification\_practices.htm
  X509v3 Authority Key Identifier:
    keyid:EB:A3:7E:89:49:47:82:DA:76:C1:AA:8F:33:65:09:A2:A7:37:BF:7C
  X509v3 Subject Key Identifier:
    D8:FB:1A:02:D0:63:59:5E:B5:BC:AC:08:96:DF:2B:34:12:42:0C:96
  X509v3 Certificate Policies:
    Policy: 1.3.6.1.4.14147.1.8.1
    CPS: http://security.fnal.gov/policies/pki\_policy\_certification\_practices.htm
  X509v3 CRL Distribution Points:
    URI:http://security.fnal.gov/pki/99f9f5a3.0
  KCA Authentication Realm:
    ..FNAL.GOV
  X509v3 Subject Alternative Name:
    email:crawdad@fnal.gov, othername:<unsupported>
```

Unsupported othername is
actually an id-pkinit-san.



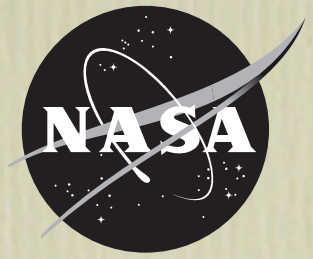
Security Issues

- All Kerberos and X.509 considerations still apply.
- Don't do PKINIT with a KX509-issued cert
 - ...unless you really know what you're doing.
 - Yes, I know loops could be fun. Might even be useful occasionally.
- Understand how Kerberos and PKI policies relate.
 - Ticket/cert lifetimes.
 - Auditing headaches getting a publicly-recognized KCA.



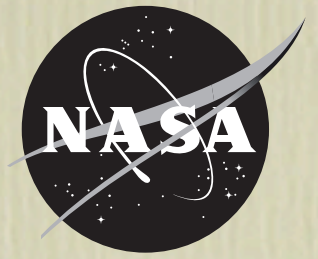
Fundamental Limitations

- Everything is in the clear.
 - Hash should protect everything's integrity.
 - Privacy/anonymity is *not* supported.
- Public key can be sniffed and reused.
 - Requestor does not have to prove knowledge of the secret key.
 - Breaks non-repudiation/digital signature applications.
 - Don't deploy with those Key Usage bits.
 - Any usage should prove knowledge of the secret key, independent of the cert.
 - TLS client OK.



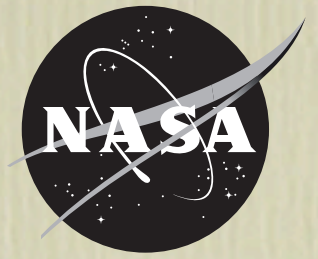
Future (Version 3.0?)

- This version has a lot of warts, mostly due to its age.
 - Originally developed with Kerberos 4!
 - Following suggestions made on IETF lists.
- Replace components of exchange with current standards which provide equivalent functionality.
 - Use Kerberos checksum instead of SHA1-HMAC
 - Probably should use KRB_XXXX or GSSAPI packaging.
 - Send PKCS-10 (RFC-2986) signed request instead of bare public key.
 - Request should also tie to Kerberos identity.
 - e-text should be UTF8, not VisibleString



Other Enhancements

- Use TCP instead of UDP
- Define a new message that says “wait”.
 - Allow for external attribute lookup operations or other complications.
- Return the entire cert chain, not just the new end-entity cert.
- Add an identifying type extension for issued certificates.



References

- Doster, W., Watts, M., and D. Hyde, "The KX509 Protocol", September 2001, <<http://www.citi.umich.edu/techreports/reports/citi-tr-01-2.pdf>>
- draft-hotz-kx509-01 <<http://tools.ietf.org/>>
 - Thanks to all the people acknowledged in Section 4 that draft.
- IETF Kerberos and PKIX mailing lists.